

Definition 0.1. A *graph* is a pair $\langle V, E \rangle$ where V is a set and E is a set of pairs from V . An *anticlique* in the graph is a set $A \subset V$ such that $[A]^2 \cap E = \emptyset$. A *chromatic number* of the graph is the smallest possible number of anticliques which together cover V . If in addition the set V is equipped with a topology, the *Borel chromatic number* of the graph is the smallest possible number of Borel anticliques which cover V .

Definition 0.2. Let $\{s_n : n \in \omega\}$ be any collection of finite binary strings such that $s_n \in 2^n$ and for every $t \in 2^{<\omega}$ there is $n \in \omega$ such that $t \subset s_n$. The graph G_0 on 2^ω is the set of all pairs of the form $\{s_n \hat{0} \hat{\ } z, s_n \hat{1} \hat{\ } z\}$ such that $n \in \omega$ and $z \in 2^\omega$.

Theorem 0.3. *The chromatic number of G_0 is 2 while the Borel chromatic number of G_0 is uncountable.*

Proof. For the evaluation of the chromatic number of G_0 , let E_0 be the equivalence relation on 2^ω connecting binary sequences $y, z \in 2^\omega$ just in case they differ on only finitely many entries. Use the axiom of choice to find a set $A \subset 2^\omega$ which visits each E_0 class in exactly one point. Let $B_0 = \{y \in 2^\omega : \text{writing } z \text{ for the unique element of } A \cap [y]_{E_0}, y \text{ and } z \text{ differ at even number of entries}\}$ and $B_1 = \{y \in 2^\omega : \text{writing } z \text{ for the unique element of } A \cap [y]_{E_0}, y \text{ and } z \text{ differ at odd number of entries}\}$. Check that B_0, B_1 are G_0 -anticliques. Clearly, $2^\omega = B_0 \cup B_1$ and so the chromatic number of G_0 is 2.

The evaluation of the Borel chromatic number uses a small claim:

Claim 0.4. *If $B \subset 2^\omega$ is a Borel non-meager set, then B contains a G_0 edge.*

Proof. The Borel set B is modulo the meager ideal equal to an open set. Since the set B is nonmeager, this means that there is a finite binary string $t \in 2^{<\omega}$ such that $B \cap [t]$ is comeager in $[t]$; let $\{O_m : m \in \omega\}$ be some countable collection of sets open dense in $[t]$ such that $\bigcap_m O_m \subset B$. Let $n \in \omega$ be such that $t \subset s_n$. By induction on $m \in \omega$ find binary strings $t_m \in 2^{<\omega}$ so that $0 = t_0 \subset t_1 \subset t_2 \subset \dots$ and both sets $[s_n \hat{0} \hat{\ } t_{m+1}]$ and $[s_n \hat{1} \hat{\ } t_{m+1}]$ are subsets of O_m . This is easily possible using the density of the set O_m . In the end, let $z = \bigcup_m t_m \in 2^\omega$ and $x = s_n \hat{0} \hat{\ } z$ and $y = s_n \hat{1} \hat{\ } z$. Both points x, y belong to $\bigcap_m O_m$ and therefore to B , and they form a G_0 -edge. \square

As a result, every Borel G_0 -anticlique is meager, and so by the Baire category theorem, countably many Borel G_0 -anticliques cannot cover 2^ω . \square

Theorem 0.5. *Suppose that G is an analytic graph on a Polish space X . Then exactly one of the following happens:*

1. *the Borel chromatic number of G is countable;*
2. *there is a continuous homomorphism of G_0 to G .*

Here, a continuous homomorphism of G_0 to G is a continuous map h from 2^ω to X such that for any $y_0, y_1 \in 2^\omega$, $y_0 G_0 y_1$ implies $h(y_0) G h(y_1)$. Note the important difference between a homomorphism and reduction: in the reduction definitions, one finds an equivalence, in the homomorphism definitions, one finds an implication.

Proof. To see that (1,2) are mutually exclusive, suppose for contradiction that $h: 2^\omega \rightarrow X$ is a continuous homomorphism of G_0 to G and $X = \bigcup_n B_n$ is a countable union of Borel G -anticliques. Then, for each $n \in \omega$, $h^{-1}B_n$ is a Borel G_0 -anticlique, and $\bigcup_n h^{-1}B_n = 2^\omega$. This, however, is impossible in view of Theorem 0.3.

To show that at least one of (1,2) must occur, start with a general preliminary claim:

Claim 0.6. *Whenever $A \subset X$ is an analytic G -anticlique, then there is a Borel anticlique $B \subset X$ such that $A \subset B$.*

Proof. Let $A_0 = \{x \in X : \exists y \in A \{x, y\} \in G\}$. The set $A_0 \subset X$ is analytic and disjoint from A , so by the Lusin separation theorem, there is a Borel set $B_0 \subset X$ containing A as a subset and disjoint from A_0 . Let $A_1 = \{x \in X : \exists y \in B_0 \{x, y\} \in G\}$. Again, this is an analytic set disjoint from A and so there is a Borel set $B_1 \subset X$ containing A as a subset and disjoint from A_1 . Let $B = B_0 \cap B_1$ and check that B works. \square

The plan now is the following: attempt to build the homomorphism as in (2) with “finite approximations”; if this attempt fails, extract from the failure countably many Borel G -anticliques covering X . We will need some notation. Fix a continuous surjection $f: \omega^\omega \rightarrow X$ and another continuous surjection $g: \omega^\omega \rightarrow G$.

Say that p is an *approximation* if $p = \langle a_p, b_p \rangle$ so that for some $n = n_p \in \omega$, a_p is a function from 2^n to ω^n and b_p is a function such that $\text{dom}(b_p) = \bigcup_{k \in n} 2^{n-k-1}$ and $\text{rng}(b_p) \subset \omega^n$.

If p, q are approximations, write $q \leq p$ if $n_p \leq n_q$, for every $t \in 2^{n_q}$ $a_q(t) \upharpoonright n_p = a_p(t \upharpoonright n_p)$, and for every $k \in n_p$ and every $t \in 2^{n_q-k-1}$, $b_q(t) \upharpoonright n_p = b_p(t \upharpoonright n_p - k - 1)$.

If $p = \langle a_p, b_p \rangle$ is an approximation, then a *validation* of p is a pair $\langle \bar{a}_p, \bar{b}_p \rangle$ such that

- $\text{dom}(\bar{a}_p) = \text{dom}(a_p)$, $\text{rng}(\bar{a}_p) \subset \omega^\omega$ and for every $t \in \text{dom}(a_p)$ $a_p(t) \subset \bar{a}_p(t)$;
- $\text{dom}(\bar{b}_p) = \text{dom}(b_p)$, $\text{rng}(\bar{b}_p) \subset \omega^\omega$ and for every $t \in \text{dom}(b_p)$ $b_p(t) \subset \bar{b}_p(t)$;
- whenever $k \in n$ is a number and $t \in 2^{n_p-k-1}$ is a binary string, then $g(\bar{b}_p(t)) = \langle f(\bar{a}_p(s_k \hat{\ } 0 \hat{\ } t)), f(\bar{a}_p(s_k \hat{\ } 1 \hat{\ } t)) \rangle$.

Claim 0.7. *Whenever $\langle p_n : n \in \omega \rangle$ is a descending sequence of approximations such that $n_{p_n} = n$ and each p_n has a validation, then there is a continuous homomorphism from G_0 to G .*

Proof. Just define the function $h: 2^\omega \rightarrow X$ by $h(y) = f(\bigcup_n a_{p_n}(y \upharpoonright n))$. It must be proved that for $y_0, y_1 \in 2^\omega$, if $y_0 G_0 y_1$ then $h(y_0) G h(y_1)$. To see this, suppose that $y_0 G_0 y_1$ holds and find $k \in \omega$ and $z \in 2^\omega$ such that $y_0 = s_k \widehat{0} z$ and $y_1 = s_k \widehat{1} z$, and let $w = \bigcup_{n > k} b_{p_n}(z \upharpoonright n - k - 1)$. It will be enough to show that $g(w) = \langle h(y_0), h(y_1) \rangle$.

if this failed, then by the continuity of the functions g, f there would have to be a number n such that for no elements $\bar{y}_0, \bar{y}_1, \bar{w}$ which agree with y_0, y_1, w respectively on the first n entries, $g(\bar{w}) = \langle h(\bar{y}_0), h(\bar{y}_1) \rangle$. This in turn means that the approximation p_n has no validation, a contradiction. \square

Thus, we are really seeking a decreasing sequence of validated approximations. The difficulty is that while p may be validated, it may still occur that none of its one step extensions are, and we have to avoid such dead ends. To do that, for a set $Y \subset X$ call an approximation p *Y-terminal* if none of its one-step extensions $q \leq p$, $q = \langle a_q, b_q \rangle$ has a validation with $\text{rng}(f \circ \bar{a}_q) \subset Y$.

Claim 0.8. *Suppose that p is an approximation and $Y \subset X$ is a Borel set. If all one step extensions of p are Y -terminal, then there are countably many Borel G -anticliques $\{B_n : n \in \omega\}$ such that p is $Y \setminus \bigcup_n B_n$ -terminal.*

Proof. Let $q \leq p$ be a one step extension of p . Write $A(q, Y) = \{f(\bar{a}_q(s_{n_q}) : \langle \bar{a}_q, \bar{b}_q \rangle\}$ is a validation of q such that $\text{rng}(f \circ \bar{a}_q) \subset Y \} \subset Y$. There are two cases.

Case 1. Suppose first that $A(q, Y)$ contains a G -edge for some $q \leq p$. Then q is not terminal: there must be two validations $\langle \bar{a}_q^0, \bar{b}_q^0 \rangle$ and $\langle \bar{a}_q^1, \bar{b}_q^1 \rangle$ of q such that $\text{rng}(f \circ \bar{a}_q^0)$ and $\text{rng}(f \circ \bar{a}_q^1)$ are both subsets of Y and $f(\bar{a}_q^0(s_{n_q}))$ is G -connected to $f(\bar{a}_q^1(s_{n_q}))$. Define the functions \bar{a}_r, \bar{b}_r as follows: $\text{dom}(\bar{a}_r) = 2^{n_q+1}$, $\bar{a}_r(t \widehat{0}) = \bar{a}_q^0(t)$, $\bar{a}_r(t \widehat{1}) = \bar{a}_q^1(t)$ for all $t \in 2^{n_q}$, $\bar{b}_r(t \widehat{0}) = \bar{b}_q^0(t)$, $\bar{b}_r(t \widehat{1}) = \bar{b}_q^1(t)$ for all $t \in 2^{n_q-k-1}$ and all $k < n_q$, and $\bar{b}_r(0) = z$ for some $z \in \omega^\omega$ such that $g(z) = \langle f(\bar{a}_q^0(s_{n_q})), f(\bar{a}_q^1(s_{n_q})) \rangle$. This latter demand can be fulfilled since by the assumptions, the $f(\bar{a}_q^0(s_{n_q}))$ is G -connected to $f(\bar{a}_q^1(s_{n_q}))$. Finally, let $r = \langle \bar{a}_r, \bar{b}_r \rangle$ be an approximation obtained from $\langle \bar{a}_q, \bar{b}_q \rangle$ by restricting all outputs of the functions \bar{a}_q, \bar{b}_q to $n_q + 1$. Then $r \leq q$ is an approximation validated by $\langle \bar{a}_r, \bar{b}_r \rangle$. This shows that q is not Y -terminal.

Case 2. Suppose that $A(q, Y)$ contains no G -edge for any one step extension $q \leq p$. Then each $A(q, Y)$ is an analytic G -anticlique, and it can be covered by a Borel G -anticlique $B(q, Y)$ by Claim 0.6. There are only countably many one-step extensions of p , and so these Borel anticliques can be enumerated by $\{B_n : n \in \omega\}$. We claim that this collection of anticliques works as in the claim.

To see this, let $Y' = Y \setminus \bigcup_n B_n$ and suppose that p is not Y' -terminal; i.e. there is a one step extension $q \leq p$ with a validation $\langle \bar{a}_q, \bar{b}_q \rangle$ such that $\text{rng}(f \circ \bar{a}_q) \subset Y'$. Now, $f(\bar{a}_q(s_{n_q})) \in A(q, Y) \subset B(q, Y)$ by the definition of $A(q, Y)$, and at the same time $\text{rng}(f \circ \bar{a}_q) \subset Y'$ and $Y' \cap B(q, Y) = \emptyset$ by the construction of the set Y' . This is a contradiction proving the claim. \square

Now, by transfinite recursion on an ordinal α , build countable sets C_α of Borel anticliques such that

- $C_0 = 0$ and $\alpha \in \beta \rightarrow C_\alpha \subset C_\beta$;
- writing $Y_\alpha = X \setminus \bigcup C_\alpha$, whenever p is an approximation all of whose extensions are Y_α -terminal, then p is $Y_{\alpha+1}$ -terminal.

This is easy to do using Claim 0.8. Note that the sets $P_\alpha = \{p: p \text{ is an approximation which is not } Y_\alpha\text{-terminal}\}$ are decreasing with α . Since the set of all approximations is countable, there must be a countable ordinal α such that $P_\alpha = P_{\alpha+1}$. There are two cases:

Case 1. The empty approximation belongs to P_α . In such a case, by induction on n build a sequence of approximations $0 = p_0 \geq p_1 \geq p_2 \geq \dots$ in P_α , using the fact that $P_\alpha = P_{\alpha+1}$ and so every approximation in P_α has a one step extension in $P_{\alpha+1}$. In the end, use Claim 0.7 to construct a continuous homomorphism h of G_0 to G .

Case 2. The empty approximation does not belong to P_α . In this case, the space X is covered by the anticliques in the set C_α . The theorem follows. \square

We will now consider, without proof, several powerful generalizations of the G_0 dichotomy.

Definition 0.9. Let $m \in \omega$ be a natural number greater than 1. Let $Y = m^\omega$ be equipped with the product topology. Let $\{s_n: n \in \omega\} \subset m^{<\omega}$ be a dense collection of strings such that $s_n \in m^n$. Let G_m be the hypergraph of dimension n on Y defined by $G_m = \{\langle s_n^s \text{mallfrown } i \hat{\ } z: i \in m \rangle: n \in \omega, z \in m^\omega\}$.

Theorem 0.10. *Let $m \in \omega$ be a natural number greater than 1.*

1. *The chromatic number of G_m is two;*
2. *the Borel chromatic number of G_m is uncountable;*
3. *whenever G is an analytic hypergraph of dimension ω on a Polish space X , then exactly one of the following occurs: either the Borel chromatic number of G is countable, or there is a continuous homomorphism $h: Y \rightarrow X$ of G_m to G .*

Definition 0.11. Let $\{s_n: n \in \omega\} \subset \omega^{<\omega}$ be a dense collection of strings such that $s_n \in m^n$. Let G_ω be the hypergraph of dimension ω on Y defined by $G_\omega = \{\langle s_n^s \text{mallfrown } i \hat{\ } z: i \in \omega \rangle: n \in \omega, z \in \omega^\omega\}$. Let $Y \subset \omega^\omega$ be the G_δ set $\{z \in \omega^\omega: \forall m \exists k > m \forall i < k \ z(i) < k\}$.

Theorem 0.12. *1. The chromatic number of G_ω is two;*

2. *the Borel chromatic number of G_ω is uncountable;*
3. *whenever G is an analytic hypergraph of dimension ω on a Polish space X , then exactly one of the following occurs: either the Borel chromatic number of G is countable, or there is a continuous homomorphism $h: Y \rightarrow X$ of G_ω to G .*

1 Applications

We will now apply the graph dichotomy theorems to prove a number of more intuitive dichotomy results. The first two applications rely on a simple observation:

Proposition 1.1. *For every $x \in 2^\omega$, the connected component $[x]$ of G_0 containing x is equal to $\{y \in 2^\omega : \{m \in \omega : x(m) \neq y(m)\} \text{ is finite}\}$.*

Proof. The left-to-right inclusion is clear since two G_0 -connected points differ at exactly one entry. For the right-to-left inclusion, for every number $n \in \omega$ consider the graph H_n on 2^n consisting of all pairs of the form $\{s_m \widehat{0} \wedge t, s_m \widehat{1} \wedge t\}$ for $m < n$ and $t \in 2^{n-m-1}$. We will prove that each graph H_n is connected.

This is proved by induction on n . The case $n = 0$ is trivial. Suppose now that H_n is connected and argue that H_{n+1} must be connected as well. Let $t_0, t_1 \in 2^{n+1}$ be arbitrary binary strings and look for a H_{n+1} -path between the two. If the last bit of t_0, t_1 is the same, then find an H_n -path from $t_0 \upharpoonright n$ to $t_1 \upharpoonright n$ and extend each string on the path by this bit. By the definitions, this yields an H_{n+1} -path from t_0 to t_1 . If the last bits of t_0, t_1 are distinct, first use the induction hypothesis to find an H_n -path from $t_0 \upharpoonright n$ to s_n and then one to s_n to $t_1 \upharpoonright n$. Then, extend the strings on the first path by the last bit of t_0 and the strings on the second path by the last bit of t_1 . The two paths taken together form an H_{n+1} -path from t_0 to t_1 as desired.

Now, if $x, y \in 2^\omega$ are points such that the set $\{m \in \omega : x(m) \neq y(m)\}$ is finite, find $n \in \omega$ larger than all elements of this set, find an H_n -path from $x \upharpoonright n$ to $y \upharpoonright n$ and append to its elements the sequence $x \upharpoonright (\omega \setminus n) = y \upharpoonright (\omega \setminus n)$. This yields a G_0 -path from x to y as desired for the right-to-left inclusion. \square

Corollary 1.2. *If $C \subset 2^\omega$ is a closed set which, with each of its elements contains also all its G_0 -neighbors, then $C = \emptyset$ or $C = 2^\omega$.*

Proof. By the previous proposition, the G_0 -components are dense. If C is nonempty, it must contain at least one whole dense G_0 -component and since C is closed, $C = 2^\omega$. \square

Theorem 1.3. (Perfect set theorem) *Let $A \subset X$ be an analytic subset of a Polish space. Then either A is countable or A contains a nonempty perfect subset.*

Proof. Consider the graph G on X connecting points x, y just in case they are distinct elements of A . The graph G is analytic, and the G_0 dichotomy gives two options.

Either, the Borel chromatic number of G is countable. Since every G -anticlique contains at most one element of A , in this case the set A must be countable. Or, there is a continuous homomorphism $h: 2^\omega \rightarrow X$ of G_0 to G . The range of h , as a continuous image of a compact space, is compact. It is also a subset of A , since every point of 2^ω does have some neighbors in G_0 and the

points in $X \setminus A$ do not have neighbors in G and the function h is a homomorphism of graphs. Thus, by the Cantor–Bendixson theorem, it will be enough to show that $\text{rng}(h)$ is uncountable. To see this, note that for every point $x \in X$ the set $h^{-1}\{x\} \subset 2^\omega$ is a closed G_0 -anticlique since h is a continuous homomorphism. Since countably many closed G_0 -anticliques cannot cover 2^ω , it follows that $\text{rng}(h)$ is uncountable as required. \square

Theorem 1.4. (Lusin–Novikov theorem) *Let $B \subset X \times Y$ be a Borel subset of a product of two Polish spaces. Then either B is the union of countably many Borel (partial) functions from X to Y , or there is a vertical section of B which contains a perfect subset.*

Proof. Let G be the graph on $X \times Y$ connecting two points just in case they are distinct elements of the same vertical section of the set B . The graph G is even Borel, and so the G_0 dichotomy gives us two options.

Either, the Borel chromatic number of G is countable. Since the intersection of every G -anticlique with the set B is a function, in this case the set B is covered by countably many functions. Or, there is a continuous homomorphism of $h: 2^\omega \rightarrow X \times Y$ of G_0 to G . First, observe that the range of h must be contained in a single vertical section of $X \times Y$. For this, note that the h -preimages of the vertical sections are closed by the continuity of h and apply Corollary 1.2. The rest of the argument is literally the same as in the perfect set theorem proof. \square

For the next two applications, we will need an old anticlique existence theorem due to Mycielski. It says that every small (meager) graph has a perfect anticlique.

Proposition 1.5. (Mycielski Theorem) *Let X be a Polish space without isolated points and G a meager graph on X . Then there is a nonempty perfect G -anticlique.*

Proof. Let O_n for $n \in \omega$ be open dense subsets of X^2 such that $G \cap \bigcap_n O_n = 0$. By induction on $n \in \omega$ build nonempty open sets P_t for all $t \in 2^n$ so that

- $t \subset s$ implies $\bar{P}_s \subset P_t$, and if t is incomparable with s then $P_s \cap P_t = 0$;
- the diameter of P_t is smaller than $2^{-|t|}$ for some fixed complete metric on X ;
- if $t \neq s$ are distinct elements of 2^n then $P_s \times P_t \subset \bigcap_{m \in n} O_m$.

In the end, define a map $h: 2^\omega \rightarrow X$ by $h(z) =$ the unique element of the intersection $\bigcap_n P_{z \upharpoonright n}$. The image of h is a perfect subset of X ; we must verify that it is an anticlique. Indeed, if $y \neq z$ are distinct elements of 2^ω and $m \in \omega$ is a number, find an $n \in \omega$ which is greater than m and such that $y \upharpoonright n \neq z \upharpoonright n$. Then, $h(y) \in P_{y \upharpoonright n}$, $h(z) \in P_{z \upharpoonright n}$, and $P_{y \upharpoonright n} \times P_{z \upharpoonright n} \subset O_m$. It follows that the pair $\langle h(y), h(z) \rangle \in X \times X$ belongs to the open set O_m for every $m \in \omega$ and so it is not in the graph G as desired. \square

The next proposition is going to be used to conclude that various graphs are meager. It ascertains that every non-meager graphon 2^ω with the Baire property contains a certain four point pattern in it. A G_0 -rectangle is a set $\{y_0, y_1\} \times \{z_0, z_1\} \subset 2^\omega \times 2^\omega$ where $y_0, y_1, z_0, z_1 \in 2^\omega$ are points such that $y_0 G_0 y_1$ and $z_0 G_0 z_1$ holds. The points y_0, y_1, z_0, z_1 (in this order) will be called the *corners* of the rectangle

Proposition 1.6. *Let G be a graph on 2^ω with the Baire property. If G is not meager then G contains a G_0 -rectangle.*

Proof. Since the graph $G \subset 2^\omega \times 2^\omega$ has the Baire property, it is equal to an open set modulo a meager set. Since G is not meager, that open set is nonempty. This means that there are binary strings $t, u \in 2^{<\omega}$ and sets O_n for $n \in \omega$ such that $O_n \subset [t] \times [u]$ is open dense in $[t] \times [u]$ and $\bigcap_n O_n \subset G$. Find incompatible binary strings t_0, u_0 extending t, u such that each of them is on the list $\{s_n : n \in \omega\}$ defining the graph G_0 . For every binary string v write $(t_0 \hat{\ } 0 \hat{\ } v)' = t_0 \hat{\ } 1 \hat{\ } v$ and similarly

Now, by induction on $n > 0$ build binary strings t_n, u_n so that

- $t_0 \hat{\ } 0 \subset t_1 \subset t_2 \subset \dots$ and $u_0 \hat{\ } 0 \subset u_1 \subset u_2 \subset \dots$
- the four open sets $[t_{n+1}] \times [u_{n+1}]$, $[t_{n+1}] \times [u'_{n+1}]$, $[t'_{n+1}] \times [u_{n+1}]$, and $[t'_{n+1}] \times [u'_{n+1}]$ are all subsets of O_n .

To perform the induction step and get t_{n+1} and u_{n+1} from t_n, u_n , just extend the binary strings t_n, u_n repeatedly four times to handle each of the four sets in the fourth item. Each time use the assumption that $O_n \subset [t] \times [u]$ is an open dense set.

In the end, let $x_0 = \bigcup_n t_n$, $x_1 = \bigcup_n t'_n$, $y_0 = \bigcup_n u_n$ and $y_1 = \bigcup_n u'_n$. By the choice of the strings t_0, u_0 , it is the case that $x_0 G_0 x_1$ and $y_0 G_0 y_1$. By the inductive construction, the G_0 -rectangle $\{x_0, x_1\} \times \{y_0, y_1\}$ is a subset of $\bigcap_n O_n$ and therefore a subset of the graph G as required. \square

Theorem 1.7. (Silver's theorem) *Let E be a coanalytic equivalence relation on a Polish space X . Then either E has countably many classes or there is a perfect set of pairwise E -unrelated elements.*

Proof. Let G be the complement of E ; this is an analytic graph. The G_0 dichotomy gives us two possibilities.

Either, the graph G has countable Borel chromatic number. In such a case, note that every G -anticlique is a subset of a single E -class, and therefore there are only countably many E -classes. Or, there is a continuous homomorphism $h: 2^\omega \rightarrow X$ of G_0 to G . The h -preimage F of E is an equivalence relation on 2^ω . It is coanalytic, and therefore has the Baire property. Note that F cannot contain a G_0 rectangle: writing y_0, y_1, z_0, z_1 for corners of such a rectangle, it would be the case that $h(y_0) E h(z_0) E h(z_1)$, by the transitivity of the equivalence relation E $h(y_0) E h(y_1)$, and this contradicts the assumption that h is a homomorphism of G_0 to the complement of E . By Proposition 1.6,

the graph F is meager, and by Proposition 1.5 there is a perfect F -anticlique $C \subset 2^\omega$, and then $h''C$ is a perfect set of pairwise E -unrelated elements. \square

For the next application, a *quasimetric* on a set X is a function $d: X^2 \rightarrow \mathbb{R}$ such that its values are non-negative, it is symmetric, and it satisfies the triangle inequality. Unlike with a metric, the set X can contain distinct points with d -distance zero. As good examples, consider the usual metric on $[0, 1]^2$ and then the metric on $[0, 1]^2$ which assigns points in the same vertical section their usual Euclidean distance, and points in distinct sections distance 1. These two metrics are not isomorphic: the former is separable, while the latter is not. The easiest way to see that the latter metric is not separable is to note that $\{(0, r) : r \in [0, 1]\}$ is an uncountable, in fact perfect, collection of points with pairwise distance ≥ 1 . The following dichotomy shows that this is the only way of proving non-separability of a definable quasimetric.

Theorem 1.8. *Let X be a Polish space and d a quasimetric on it such that for every $\varepsilon > 0$, the set $d^{-1}[\varepsilon, \infty) \subset X \times X$ is analytic. Then either d is separable or there is $\varepsilon > 0$ and a perfect set of elements of X with pairwise distances greater than ε .*

Proof. For each $\varepsilon > 0$ let G_ε be the graph on X connecting x, y if $d(x, y) \geq \varepsilon$; this is an analytic graph. The G_0 dichotomy gives us two options.

Either, for every $\varepsilon > 0$, the Borel chromatic number of G_ε is countable. Note that every G_ε -anticlique is a set of points of pairwise distance $< \varepsilon$, and so in this case for every $\varepsilon > 0$ the space X can be covered by countably many ε -balls, for every $\varepsilon > 0$. This implies that the metric d is separable.

Or, there is a real $\varepsilon > 0$ and a continuous homomorphism $h: 2^\omega \rightarrow X$ of G_0 to G_ε . Let e be the quasimetric on 2^ω defined by $e(y, z) = d(h(y), h(z))$ and consider the graph H on 2^ω connecting points y, z if their e -distance is $\leq \varepsilon/2$. The graph H is coanalytic and therefore has the Baire property. It cannot contain a G_0 -rectangle: if y_0, y_1, z_0, z_1 were corners of such a rectangle, then $d(h(y_0), h(z_0)) \leq \varepsilon/2$, $d(h(y_1), h(z_0)) \leq \varepsilon/2$ would hold. At the same time, $d(h(y_0), h(y_1)) > \varepsilon$ since h is a homomorphism of G_0 to G_ε . This contradicts the triangle inequality for the quasimetric d . By Proposition 1.6, the graph H is meager, and by Proposition 1.5 there is a perfect H -anticlique $C \subset 2^\omega$. Then $h''C$ is a perfect set of points of pairwise d -distance at least $\varepsilon/2$. This completes the proof. \square

The next application concerns linear quasiorders on Polish spaces. Here, a *quasiorder* on a set is a transitive relation \leq which contains the diagonal. It is *linear* if for all x, y in its domain, either $x \leq y$ or $y \leq x$. To motivate the result, consider the linear ordering \leq on \mathbb{R} and the linear ordering \prec on \mathbb{R}^2 defined by $\langle x_0, x_1 \rangle \prec \langle y_0, y_1 \rangle$ if $x_0 < x_1$ or $x_0 = x_1$ and $y_0 \leq y_1$. They are not isomorphic since \leq has a countable dense set while \prec does not. One way to see that \prec does not have a countable dense set is to observe that $[\langle x, 0 \rangle, \langle x, 1 \rangle]$ for $x \in \mathbb{R}$ is an uncountable collection of pairwise disjoint intervals in \prec . The following dichotomy shows that this is in fact the only way of showing that a definable linear quasiorder is not separable.

Theorem 1.9. *Let X be a Polish space and \leq an analytic linear quasiordering on X . Then either \leq is separable or else there is a perfect collection of pairwise disjoint \leq -intervals.*

Proof. We will need a notational convention. Say that a pair $\langle x, y \rangle \in X^2$ determines a nontrivial interval in \leq if either $x \leq y$ or $y \leq x$ fails; the closed interval determined by the pair will be the set $\{z \in X : x \leq z \leq y\}$ or $\{z \in X : y \leq z \leq x\}$ depending on whether $\leq y$ or $y \leq x$ fails. We will identify the pairs with the closed intervals that they determine.

Consider the graph G on $X \times X$ connecting pairs $\langle x_0, x_1 \rangle$ and $\langle y_0, y_1 \rangle$ if they determine nontrivial closed \leq -intervals and these intervals have empty intersection. An inspection reveals that this is an analytic graph. The G_0 dichotomy provides us with two options.

Either, the (Borel) chromatic number of the graph G is countable. In such a case, the nontrivial closed intervals can be organized into countably many sets B_n for $n \in \omega$ such that in each B_n , any two intervals have nonempty intersection. In this case, we will argue that the completion \leq^* of the quasiordering \leq is separable, which certainly means that \leq is separable. For each interval $I \in B_n$ write I^* for the corresponding interval in \leq^* . Let x_n^* be the least upper bound of the left endpoints of intervals in B_n , and observe that $x_n^* \in \bigcap \{I^* : I \in B_n\}$. The set $\{x_n^* : n \in \omega\}$ is dense in \leq^* since every nontrivial closed interval in \leq^* contains one of the points x_n^* .

Or, there is a continuous map $h : 2^\omega \rightarrow X \times X$ which is a homomorphism of G_0 to G . Note that for every $z \in 2^\omega$ the pair $h(z)$ defines a nontrivial interval in \leq , since only such pairs have any neighbors in the graph G . We must find a perfect set $C \subset 2^\omega$ such that $h''C$ consists of pairs defining pairwise disjoint intervals. Let H be the graph on 2^ω which is the complement of the h -preimage of G . This is a coanalytic graph, and therefore has the Baire property. It cannot contain a G_0 -rectangle; to see this, observe that if I, J are disjoint closed intervals and K, L are closed intervals intersecting both I, J in a nonempty set then K, L cannot be disjoint. Thus, the graph H is meager by Proposition 1.6, and so by Proposition 1.5, there is a perfect H -anticlique $C \subset 2^\omega$. Then, $h''C$ is the requested perfect collection of pairwise disjoint intervals. \square